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A 0.3-V Operating, V_{th} -Variation-Tolerant SRAM under DVS Environment for Memory-Rich SoC in 90-nm Technology Era and Beyond

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Summary

We propose a voltage control scheme for 6T SRAM cells that makes a minimum operation voltage down to 0.3 V under DVS environment. A supply voltage to the memory cells and wordline drivers, bitline voltage, and body bias voltage of load pMOSFETs are controlled according to read and write operations, which secures operation margins even at a low operation voltage. A self-aligned timing control with a dummy wordline and its feedback is also introduced to guarantee stable operation in a wide range of the supply voltage. A measurement result of a 64-kb SRAM in a 90-nm process technology shows that a power reduction of 30% can be achieved at 100MHz. In a 65-nm 64-Mb SRAM, a 74% power saving is expected at 1/6 of the maximum operating frequency. The performance penalty by the proposed scheme is less than 1%, and area overhead is 5.6%.

Key words: SRAM, DVS, V_{th} -variation-tolerant, low power

1. Introduction

In order to save a power of an SoC, dynamic voltage scaling (DVS) that adaptively controls an operating frequency and supply voltage (V_{dd}) has been implemented in a mobile system [1]. However, a minimum operation voltage (V_{min}) is becoming higher as a fabrication technology is scaled down, since operation margins of memory cells in an embedded SRAM are degraded under both read and write conditions due to threshold-voltage (V_{th}) variation of MOSFETs. Figure 1 illustrates so-called Pelgrom plots [2], which shows characteristics of standard deviation in 90-nm and 65-nm CMOS process technologies. A standard deviation of V_{th} ($\sigma_{V_{th}}$) has been

formulated as follows [3].

$$\sigma_{V_{th}} \propto T_{OX} \cdot \frac{\sqrt[4]{NT \cdot \ln(N/n_i)}}{\sqrt{L_{eff}W_{eff}}}, \quad (1)$$

where T_{OX} is a gate oxide thickness, N is a channel dopant concentration, T is a absolute temperature, n_i is a intrinsic carrier concentration, and L_{eff} and W_{eff} are an effective channel length and width of MOSFETs, respectively. The 65-nm process technology has the smaller slopes in the Pelgrom plots thanks to thinner T_{OX} . However, the V_{th} variation in the 65-nm technology is larger than that in the 90-nm one because of its smaller $L_{eff}W_{eff}$. Since a chip area of 80% or more is supposed to be occupied with memories [4] and unfortunately a V_{th} deviation is getting larger, V_{min} will be restricted by the V_{th} variation, which hinders wide-range power scaling of a future SoC with DVS capability.

Fig. 1. Pelgrom plots in 90-nm and 65-nm process technologies.

Some techniques have been proposed to improve operation margins of memory cells. One is a power-line-floating in write operation for the conventional 6T memory cells [5], and another is a 7T cell which secures read margin by cutting a loop of a memory-cell flip-flop in read operation [6]. However, these techniques only improve either read or write margin, and thus V_{min} is still limited by the other operation. Both read and write margins are

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important. The voltage switching using the dual-Vcc scheme is also proposed for improvement in both read and write margins [7], but it is not sufficient for DVS.

In this paper, we report an optimum voltage control scheme for 6T memory cells that lower V_{min} under the DVS environment. This scheme selectively controls voltages in the memory cell to expand the read and write margins. A self-aligned timing control is also developed so as to secure a timing sequence of signals against process, voltage and temperature (PVT) variation.

The rest of this paper is organized as follows. Section 2 introduces the proposed optimum voltage control scheme, and shows the improvement of the operation margins in simulation. The circuit implementation of the proposed scheme and self-aligned timing control are also discussed in Section 2. In Section 3, experimental results of a 64-kb SRAM in a 90-nm CMOS technology are demonstrated, including a measured fail-bit count of memory cells and relationship between power and operating frequency. Simulation results in a case that a memory capacity increases and process technology is scaled down to 65 nm, are also shown in Section 3. Conclusions are given in Section 4.

2. Optimum Voltage Control Scheme

2.1 Concept

Fig. 2. Block diagram of SoC with DVS.

Under the proposed DVS environment, a high supply voltage (V_{max}) is applied externally, and a variable supply voltage (V_a) from a DC/DC converter is also provided to logic and SRAM modules in an SoC, as shown in Fig. 2. V_a is adaptively controlled, which is between V_{min} and V_{max} . In this paper, V_{max} is set to 1.0 V as a nominal voltage in a 90-nm process technology. In the optimum voltage control scheme proposed, both V_a and V_{max} are provided to the SRAM, and a supply voltage, wordline (WL) and bitline (BL) voltages of the memory cells are switched according to the read and write conditions in

Table 1 and Fig. 3. Figure 3 is schematics of a 6T memory cell, where V_{mc} , V_{wl} , and V_{bl} are a supply voltage, a wordline voltage, and a bitline voltage of a cell under the read and write conditions, respectively, and V_1 and V_2 are voltages of data-stored nodes.

Table 1. Voltage controls in the conventional and proposed SRAM.

Part \ State	Conv.	Proposed		
		Read	Write	Non-access
Peripheral Vdd	V_a	V_a	V_a	V_a
BL pre-charge level	V_a	V_a	V_a	V_a
WL swing	V_a	V_a	V_{max}	-
MC Vdd	V_a	V_{max}	V_a	V_a
MC Vbp	V_a	V_{max}	V_{max}	V_{max}

V_{max} : Maximum supply voltage
 V_a : Variable supply voltage under DVS environment
 BL: Bitline
 WL: Wordline
 MC: Memory cell

Fig. 3. Schematics of a 6T memory cell, the worst-case conditions of local V_{th} variations, and voltage controls in the proposed scheme. (a) Read condition, and (b) write condition.

The optimum voltage control scheme improves operation margins in both read and write operations. A supply voltage of memory cells is set to V_{max} in the read operation to maximize the read margin. Alternatively in a write cycle, a WL voltage is set to V_{max} to obtain the write margin. An n-well bias (V_{bp}) for load pMOSFETs in the memory cells is tied to V_{max} , which increases the

pMOSFET V_{th} and write margin when V_a is less than V_{max} . Although the negative V_{bp} is utilized, a serious problem of a gate-induced drain leakage (GIDL) or negative bias temperature instability (NBTI) does not occur since V_{bp} does not exceed V_{max} . Moreover, this scheme can be implemented with a twin-tub process technology, and soft-error immunity is higher than that in a triple-well technology. The improvement of the operation margins with the proposed voltage control scheme is given in the next subsection.

2.2 Improvement of Operation Margins

Fig. 4. (a) Read margins and (b) write margins at $n=0$ (no V_{th} variation) and $n=3$.

In order to estimate operation margins at each process corner, the worst case of local V_{th} variations (random variations) of MOSFETs in a memory cell are set as indicated in Fig. 3. $\sigma_{V_{th}}$ is a standard deviation of V_{th} in a MOSFET, but strictly the values are different among process corners, which will be discussed later. n is a coefficient. Four transistors in a memory cell (AC1, DR1, DR2 and LD2 for read; AC1, LD1, DR2 and LD2 for write in Fig. 3) affect operation margins [8][9], and $n=3$, for example, indicates that a local V_{th} variation of $6\sigma_{V_{th}}$ is considered in a memory cell. The read margin is determined by a logical V_{th} of an inverter in a memory-cell flip-flops (LD2 and DR2), and a voltage of a data-stored node, V_1 , when “L” is stored in V_1 and a WL is activated. The worst case of the local V_{th} fluctuation for the read margin is that the logical V_{th} of the inverter is lowest, and that the ratio of the conductances between DR1 and AC1 are smallest. As for the write margin, the logical V_{th} of an inverter and the voltage of V_1 are important, when V_1 is already “H” and then “L” is being written. The write margin is getting worse at a lower logical V_{th} of the inverter and a larger ratio of the conductances between LD1 and AC1. Figure 4 (a) is

called a butterfly plot and illustrates a read margin in a memory cell at $n=0$ (no V_{th} variations) and $n=3$. Figure 4 (b) shows a definition of a write margin [10]. The read and write margins worsen if V_{th} variation occurs, which raises V_{min} .

Figure 5 is a so-called milky-way plot of a memory cell, and demonstrates the operation margins at three V_{th} points (A-C). A diamond illustrated at the center in Fig. 5 shows process corners (FF, FS, SF, SS, and CC). Global V_{th} variation (wafer-to-wafer / lot-to-lot variation) is reflected to the process corners. The local V_{th} variation of $6\sigma_{V_{th}}$ described above is also considered in this figure. Point (A) is out of the read limit, and there is no read margin as shown in the bottom left butterfly curves. Point (A) explains the situation that read margin cannot be obtained and V_1 of “L” will not be kept in during a read operation, which is due to the lower logical V_{th} of the inverter. Similarly Point (B) is out of the write limit and write operation will fail, since the voltage of V_1 is much higher than 0 V when “L” is written, which is caused by the larger conductance ratio between LD1 and AC1. Both read and write operations will pass at Point (C) where the logical V_{th} of the inverter is adequately high and the conductance ratio between LD1 and AC1 is much smaller. In the region between the read and write limit lines, a memory cell properly works.

Fig. 5. Milky-way plot of a memory cell when $V_{dd}=1.0$ V. At Point (A), there is no read margin while there is at Point (C). Write margin exists at Point (C) while it does not at Point (B).

Fig. 6. Milky-way plots of (a) the conventional memory cell, and (b) the proposed one using optimum voltage control scheme.

Fig. 7. Block diagram of the proposed 64-kb SRAM with the optimum voltage control scheme. The control signals, WE, SE, and PC_n are input to write circuit, sense amplifier, and bitline pre-charger, respectively. These signals are also utilized as control signals for other logic gates.

Figures 6 (a) and (b) show the detailed milky-way plots in the conventional and proposed optimum voltage control schemes. In the conventional SRAM, the margins become smaller as V_a is decreased. In particular at a V_a of 0.6 V, there is neither read nor write margin. To the contrary, the proposed scheme ensures sufficient read and write margins even if V_a is low, and the margins rather becomes larger as V_a is reduced. The expansion of read margin is due to a lower voltage of V_1 when V_1 is “L”. The gate voltages of DR1 and AC1 are V_{max} and V_a , respectively, hence the conductance ratio between DR1 and AC1 is larger. Similarly in the write operation of “L”, the voltage of V_1 is sufficiently low since the conductance ratio between LD1 and AC1 is small, which is attributed to the gate voltage of AC1 (V_{max}). This is why the write margin is widening when V_a is low.

2.3 Circuit Design

Figure 7 illustrates a block diagram of the proposed SRAM to which the optimum voltage control scheme is applied. In the proposed SRAM, a memory cell array is divided into 64 blocks so that one block has 128 words by 8 bits, in which the voltage controls are done by a block-by-block basis since Vdd lines in the memory cells are along with bitlines (BLs) [11] unlike the row-by-row Vdd

control [12]. Vdd selectors are implemented in order to change the voltages (V_{mc} and WL voltage). A Vdd selector consists of switches of p-channel MOSFETs, and the total channel width of the pMOSFET switches is set to 20-fold of the total channel width of MOSFETs in a memory cell. This value minimizes the Vdd switching time. If the switch width is small, the Vdd switching time becomes longer. On the other hand, even if the switch width is large, it still takes longer time to drive the gate of the switch.

Level shifters are introduced just after X decoders to amplify the WL voltages to V_{max} in a write cycle. The output voltage of the Vdd selector for WLs (V_{wl}) is provided to an AND gate of GWL and Clwl, which controls local WL (LWL) voltages. Figure 8 demonstrates the read margins in a write cycle when V_a is fixed to 0.4 V and a voltage on a WL is varied. The margin could not be obtained if the voltage of WL was V_{max} (1.0 V). Again, this means that stored data would be destroyed under the write condition, since stored data in other blocks are destroyed if the conventional single WL scheme is used. Therefore, the divided-WL structure [13] that can hierarchically access to a local WL is applied in the proposed SRAM, in order to make a WL voltage in non-accessed memory cells grounded. Besides, a channel length of access MOSFETs (AC1 and AC2 in Fig. 3) is set

to $0.15\ \mu\text{m}$, which is longer than the minimum rule of $0.10\ \mu\text{m}$. This does not only improve the read margin but also reduces a BL leakage current within a marginal delay overhead. The delay overhead is discussed in Section 3.3.

2.4 Self-Aligned Timing Control

A proper sequence of the voltage controls is of importance to avoid unexpected flips in memory cells. As pointed out in the previous subsection, when a cycle is changing from a write (or non access) to read operation, a destructive read might occur because a WL voltage gets higher than V_{mc} . In order to clear this issue, a timing adjustment between the WL voltage and V_{mc} in Fig. 9 (a) is necessary. Moreover as shown in Fig. 9 (b), another sequence that a WE signal is negated after a WL voltage is grounded, is required as well when a write operation is concluded. Otherwise, the bitline voltages would be floating while the WL voltage is high, and wrong data would be written to memory cells.

Fig. 8. Butterfly plots at voltage when $V_a=0.4\text{ V}$ and a WL voltage is varied.

Fig. 9. Timing chart of voltage controls.

In order to secure these sequences, a self-aligned timing control is implemented with a dummy WL and its feedback (see Fig. 7). C_{vmc} in Fig. 7 is a control signal for memory cell power supply (V_{mc}), and V_{dd} selector for memory cells chooses V_{max} if C_{vmc} is grounded.

C_{vwl} means a control signal for a local WL voltage (V_{wl}), and V_{wl} is raised to V_{max} when C_{vwl} is pulled down. At the beginning of a read cycle and when a row and a column are selected, a control signal of RS-Read and CS are pulled up and C_{vmc} is grounded, which results in selecting V_{max} for V_{mc} . Then a local WL (LWL) is pulled up to V_a after C_{vmc} and global WL (GWL) are activated, which secures the timing sequence in Fig. 9 (a). We confirmed that, in a simulation, the time required for the voltage switching is negligible compared to the cycle time because of a small capacitance in a memory-cell block power supply. In contrast, after data are written to cells, clock signal is pulled down and the dummy WL starts falling. A WE signal is deactivated after receiving the feedback signal of the dummy WL. The use of the dummy WL guarantees the timing sequences in Fig. 9 (b) against PVT fluctuation.

3. Simulation and Measurement Results

3.1 Chip Implementation

A 64-kb SRAM test chip was designed and fabricated in a 90-nm CMOS process technology to verify the feasibility of the proposed scheme. Figure 10 shows a micrograph of the test chip and a layout view of a memory-cell block. The area overhead of the proposed SRAM is 5.6%, which is basically caused by the V_{dd} selectors and level shifters.

Fig. 10. Chip micrograph and layout view of a memory-cell block.

3.2 Fail-Bit Count

Figure 11 (a) demonstrates the proposed scheme improving the operation margins, by means of a measured fail-bit count (FBC) in a read operation. V_{min} is limited by read operation in our design, since the process corners

are closer to the read limit curve in the milky-way plot as indicated in Fig. 6 (a). The clock-cycle time in this measurement is as slow as 1 μ s in order to evaluate $V_{\min}$, which is 0.55 V in the conventional SRAM while that of the proposed one is as low as 0.3 V at the CC process corner. In the conventional scheme, the V_{th} variation of the memory cells governs $V_{\min}$. However in the proposed scheme, the value of 0.3 V means a lower limit of peripheral circuits since the memory cells have a larger margin as V_a is decreased as already shown in Fig. 6 (b). In the proposed SRAM, the GWL level shifter does not work below 0.3V.

Note that in the conventional read/write operation, $V_{\min}$ gets higher at the FS corner as shown in Fig. 11 (b) because the FS corner is closest to the read limit line in the milky-way plot. In addition, a bit error rate in data retention should be considered because the transistor V_{th} of cross-coupled inverters in a memory cell gets unbalanced as the V_{th} variation increases. Thus, the stored data cannot be kept even if memory cells are not accessed. The simulated curves in the conventional read/write operation and data retention are obtained as follows:

- Minimum operation voltages of memory cells are estimated with a bit error rate, which is defined as the inverse of the capacity. Coefficient n in Fig. 3 is set to 2.17 for 64-kb, 2.52 for 2-Mb, and 2.83 for 64-Mb SRAMs.
- At FS corner, the read margin is degraded further since $\sigma_{V_{th}}$ of load pMOSFETs, as expressed in (1), is larger due to a higher channel dopant concentration, N . This leads to get a BER- V_a curve more gradual and raises $V_{\min}$.

Fig. 11. (a) Measured FBC, and (b) simulated BER when a process corner and memory capacity are varied.

If a memory capacity is assumed to be 64 Mb, the conventional scheme does not work below 0.79 V in a 90-nm process technology, which hinders the DVS

advantages. On the contrary, $V_{\min}$ in the proposed scheme in a 64-Mb SRAM can be down to 0.36 V which is restricted by data-retention characteristics. BER in a 65-nm process technology is also shown in Fig. 11 (b). Since local V_{th} variation is larger as a process technology is scaled down as illustrated in Fig. 1, $V_{\min}$ becomes as high as 0.95 V in the conventional 65-nm 64-Mb SRAM, which means that there is just a 0.05-V room to change the supply voltage, and thus DVS never function under that condition.

3.3 Power-versus-Frequency Characteristics

Figure 12 (a) shows power dependences on an operating frequency (P-f curves) in the designed 90-nm 64-kb SRAM. V_a is implicitly adjusted according to the operating frequency. The performance penalty by applying the proposed scheme is less than 1% when V_a is 1.0 V. V_a must not be lowered than 0.66 V in the conventional scheme, which results in a higher power at a frequency of less than 300 MHz. The proposed scheme in the measurement allows the lower-voltage operation at the lower frequency, and we confirmed by the measurement that the 100-MHz 0.45-V operation has an advantage of 30% power reduction compared with the conventional scheme. As indicated in Fig. 12 (b), power savings are higher as a memory capacity increases and process technology is scaled down since $V_{\min}$ in such situation becomes higher. A power saving of 57% in a 90-nm 64-Mb SRAM and 74% in a 65-nm 64-Mb SRAM can be achieved at 1/6 of the maximum operating frequency.

Fig. 12. P-f curves in (a) 90-nm 64-kb SRAM and (b) various capacities and process technologies.

4. Summary

This paper presented an optimum voltage control scheme in memory cells and self-aligned timing control method under DVS environment which expands read and write

operation margins and allows an operation at as low as 0.3 V in a 90-nm SRAM. The proposed scheme achieves a power saving of 30% at 100MHz in a 90-nm 64-kb SRAM, and power reduction of 74% at 1/6 of the maximum frequency in a 65-nm 64-Mb SRAM. Since the proposed scheme makes V_{min} much lower as memory capacity increases and process technology is scaled down, DVS can be enjoyed even on a future memory-rich SoC.

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